HighLoad++

Performance Enhancements in PostgreSQL 8.4

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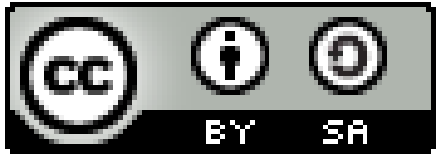

PostgreSQL 8.4

- Released July 2009
  - 8.4.1 released September 2009
- Major upgrade from 8.3
- New features and enhancements of existing ones

Using PostgreSQL performance

- “ORM-like queries” only get you so far
- Application specific optimizations
- Don't be afraid to let the database work!

Performance enhancements

- Some are application transparent
  - Possibly even DBA transparent
- Some require application changes

Let's get started

- Query execution optimizations

Anti-joins and Semi-joins

- Formalized JOIN methods for inequality joins
- Better performance for EXISTS / NOT EXISTS

Anti-joins and Semi-joins

- 8.3

```
pagila=# EXPLAIN SELECT * FROM actor a WHERE NOT EXISTS
   (SELECT * FROM film_actor fa WHERE fa.actor_id=a.actor_id);

 Seq Scan on actor  (cost=0.00..288.99 rows=100 width=25)
   Filter: (NOT (subplan))
   SubPlan
     ->  Index Scan using film_actor_pkey on film_actor
           (cost=0.00..38.47 rows=27 width=12)
         Index Cond: (actor_id = $0)
```

Anti-joins and Semi-joins

- 8.3

```
pagila=# EXPLAIN SELECT * FROM actor a WHERE NOT EXISTS
   (SELECT * FROM film_actor fa WHERE fa.actor_id=a.actor_id);

 Nested Loop Anti Join  (cost=0.00..30.57 rows=1 width=25)
   ->  Seq Scan on actor  (cost=0.00..4.00 rows=200 width=25)
   ->  Index Scan using film_actor_pkey on film_actor
         (cost=0.00..1.54 rows=27 width=2)
         Index Cond: (film_actor.actor_id = actor.actor_id)
```

Anti-joins and Semi-joins

- 8.3

```
pagila=# EXPLAIN SELECT * FROM actor a WHERE EXISTS
   (SELECT * FROM film_actor fa WHERE fa.actor_id=a.actor_id);

 Nested Loop Semi Join  (cost=0.00..30.57 rows=200 width=25)
   ->  Seq Scan on actor  (cost=0.00..4.00 rows=200 width=25)
   ->  Index Scan using film_actor_pkey on film_actor
         (cost=0.00..1.54 rows=27 width=2)
         Index Cond: (film_actor.actor_id = actor.actor_id)
```

Hash for DISTINCT/UNION

- Previously, always a sort+unique
- *No longer guaranteed sorted!*
  - Add ORDER BY
  - Both plans will be considered
- Also affects EXCEPT & INTERSECT

Hash improvements

- Faster algorithms
- Also faster hash indexes
  - Still not WAL-logged
- And optimizations of HASH joins
  - Particularly around large joins

Moving on

- DBA optimizations

Function level statistics

- pg_stat_user_functions
- Controlled by “track_functions”
  - *none, pl or all*
- Tracks calls, time, and internal time

```
postgres=# select * from pg_stat_user_functions ;
-[ RECORD 1 ]------
funcid      | 101414
schemaname  | public
funcname    | foo
calls       | 1003
total_time  | 6
self_time   | 6
```

Free Space Map (FSM)

- Stores list of free blocks in relations
  - Caused by DELETE and UPDATE
- Used by INSERT & UPDATE

New Free Space Map (FSM)

- No more max_fsm_pages!
- Dynamically tuned
- Uses normal buffer cache

New Free Space Map (FSM)

- No global lock
- Not lost on crash

New Free Space Map (FSM)

- No global lock
- Not lost on crash
- VACUUM is still needed, of course...

Visibility Map

- Tracks pages that are “visible to all transactions” in bitmap
- Set by VACUUM
- Cleared by INSERT/UPDATE/DELETE

Partial VACUUM

- “Visible to all” pages skipped by VACUUM
- Only heap tables, not indexes
- Still requires freezing

VACUUM snapshot tracking

- Snapshot tracking for idle sessions
- Makes VACUUM clean up better with long running transactions
- <IDLE> In Transaction

Stats temp file improvements

- Previously, unconditionally written twice/sec in data dir
- Now, written only on demand
- And in configurable location (tmpfs!)

Parallel pg_restore

- Restore from dump was single threaded
- Can now load in <n> sessions
- At least one table per session
- No single-transaction!

int8 pass by value

- 64-bit integers finally take advantage of 64-bit CPUs

Moving on

- Application features

Subselects in LIMIT/OFFSET

- Previously, only constants allowed
- Required two queries / roundtrips
  - Or cursor in function
- **`SELECT * FROM … LIMIT (`**

```
    SELECT something FROM other
)
```

WINDOW aggregates

- Perform aggregates over parts of data
- Avoid requiring multiple queries
- Avoid multiple scans

```
SELECT name, department, salary,
  rank() OVER (
    PARTITION BY department
    ORDER BY salary DESC
  )
FROM employees
```

name	department	salary	rank
Berra	Ekonomi	29400	1
Åke	Ekonomi	29400	1
Sune	Ekonomi	24000	3
Arne	IT	24000	1
Pelle	IT	22000	2
Kalle	IT	18000	3

(6 rows)

```
SELECT name, department, salary,
  rank() OVER (
    PARTITION BY department
    ORDER BY salary DESC
  ),
  rank() OVER (
    ORDER BY salary DESC)
FROM employees
```

Common Table Expressions

- WITH RECURSIVE
- Traverse trees and graphs in SQL
- .. avoid multiple queries
  - (also makes your life easier)

```
WITH RECURSIVE t(id, department, name, manager) AS
(
  SELECT id, department, name, manager
   FROM emp WHERE name='Kalle'
 UNION ALL
  SELECT emp.id,emp.department,emp.name,emp.manager
   FROM emp JOIN t ON t.manager=emp.id
)
SELECT * FROM t;
```

```
 id | department  | name   | manager
----+-------------+--------+---------
  1 | IT          | Kalle  |       3
  3 | IT          | Arne   |       5
  5 | Ekonomi     | Berra  |
(3 rows)
```

```
 id | department  | name   | manager
----+-------------+--------+---------
  1 | IT          | Kalle  |       3
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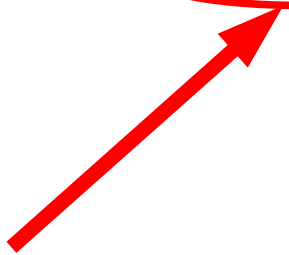

Very important!

Lots of more improvements!

- But that's it for now..
- Go download and test!

HighLoad++

Questions?

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